

ACM International Collegiate Programming Contest — Training Session III

C. Maria Keet

Department of Computer Science, University of Cape Town, South Africa
mkeet@cs.uct.ac.za

August 27, 2016

Today

- Graphs (other than just plain Dijkstra)
- Some more and less common data structures
- Practice with the Arab Regional 2013

Outline

- 1 Graphs
 - As long as I learn, I live
 - Railway construction
 - Software Allocation
- 2 Data structures
 - Streaming System
 - Young and Successful
 - Hic-Hac-Hoe
- 3 Strategy and teamwork

As long as I learn, I live

- UVa 459 Graph connectivity / 12376 As long as I learn, I live
- Graphs

As long as I learn, I live

- UVa 459 Graph connectivity / 12376 As long as I learn, I live
- Graphs
- ⇒ Finding connected components type of problem, using DFS or BFS
- ⇒ See also §4.2.3 of CP3

Designing Railway routes

- 2013 regionals problem (see printout), few solved it
- type: algorithmically, somewhat challenging
- Solve it

Designing Railway routes

- 2013 regionals problem (see printout), few solved it
 - type: algorithmically, somewhat challenging
 - Solve it
- ⇒ My direction of solution is one of 'modify existing algorithm'

Designing Railway routes: toward solution

- Task: second-best solution
- Has lots of constraints on the input data
- Check first example, and draw and check the second example (in the input/output files, given)



Designing Railway routes: toward solution

- Task: second-best solution
- Has lots of constraints on the input data
- Check first example, and draw and check the second example (in the input/output files, given)
- Do you have to start at node 1?



Designing Railway routes: toward solution

- Task: second-best solution
- Has lots of constraints on the input data
- Check first example, and draw and check the second example (in the input/output files, given)
- Do you have to start at node 1? no.
- What algorithms do we have to compute trees and shortest paths? Can we use that here?



Designing Railway routes: toward solution

- Task: second-best solution
- Has lots of constraints on the input data
- Check first example, and draw and check the second example (in the input/output files, given)
- Do you have to start at node 1? no.
- What algorithms do we have to compute trees and shortest paths? Can we use that here?
- OSPF/Dijkstra's algorithm, spanning tree algorithms



Designing Railway routes: toward solution

- Task: second-best solution
- Has lots of constraints on the input data
- Check first example, and draw and check the second example (in the input/output files, given)
- Do you have to start at node 1? no.
- What algorithms do we have to compute trees and shortest paths? Can we use that here?
- OSPF/Dijkstra's algorithm, spanning tree algorithms
- But can we get the 2nd-best from that?
 - Compute OSPF for each node, take the 2nd-lowest value, or mutate lowest value?
 - Modify a greedy algorithm to find 'second-best'?
 - Which one works best? or neither, and use another solution?

Another step toward solution

- Dijkstra (or other link state algorithm): no (why not?)
- Bellman-Ford (distance vector algorithm)? no (same reason as previous)

Another step toward solution

- Dijkstra (or other link state algorithm): no (why not?)
- Bellman-Ford (distance vector algorithm)? no (same reason as previous)
- need a **minimum spanning tree**, and take the 2nd best option
- Or perhaps compute all, rank them on cost, take the 2nd best?

Another step toward solution

- Dijkstra (or other link state algorithm): no (why not?)
 - Bellman-Ford (distance vector algorithm)? no (same reason as previous)
 - need a **minimum spanning tree**, and take the 2nd best option
 - Or perhaps compute all, rank them on cost, take the 2nd best?
 - What is a minimum spanning tree?
-
- Which algorithms do we have?

Another step toward solution

- Dijkstra (or other link state algorithm): no (why not?)
- Bellman-Ford (distance vector algorithm)? no (same reason as previous)
- need a **minimum spanning tree**, and take the 2nd best option
- Or perhaps compute all, rank them on cost, take the 2nd best?
- What is a minimum spanning tree?
 - A spanning tree of a connected, undirected, graph is a subgraph that is a tree and connects all the vertices together.
- Which algorithms do we have?

Another step toward solution

- Dijkstra (or other link state algorithm): no (why not?)
- Bellman-Ford (distance vector algorithm)? no (same reason as previous)
- need a **minimum spanning tree**, and take the 2nd best option
- Or perhaps compute all, rank them on cost, take the 2nd best?
- What is a minimum spanning tree?
 - A spanning tree of a connected, undirected, graph is a subgraph that is a tree and connects all the vertices together.
- Which algorithms do we have?
 - Skiena (pp192-202): Prim's algorithm (also greedy), Kruskal's (also greedy, but doesn't start with a particular vertex)
 - Longer list: https://en.wikipedia.org/wiki/Minimum_spanning_tree#Algorithms

Solution (sketch)

1. Kruskal for MST
<https://www.youtube.com/watch?v=71UQH7Pr9kU>
2. For each of the non-MST edges:
 3. Take edge not in MST that results in a cycle, ϵ
 4. Compare weight ϵ with heaviest (non- ϵ) edge in the cycle
 5. If lowest value so far, then store as lowest solution σ
6. add ϵ of best solution σ and remove the heaviest (non- ϵ) edge of that cycle

Software Allocation

- UVa 259 Software Allocation
- Graphs
- Note: this needs a different class of graph algorithms cf. what we have seen so far
- Figure out how your graph look like,
- What is different about it cf. the others?

Software Allocation

- UVa 259 Software Allocation
 - Graphs
 - Note: this needs a different class of graph algorithms cf. what we have seen so far
 - Figure out how your graph look like,
 - What is different about it cf. the others?
- ⇒ You need a 'network flow' algorithm for this

What is a 'network flow'?

- Directed graph, where each edge has a max capacity
- Amount of 'flow' into the node = amount of 'flow' out of the node, unless it's the source or sink
- Used for modelling traffic systems, fluids in pipes, bits across wires...
- Can we indeed model the Software Allocation problem as a network flow one? Show it.
- Then: which algorithm to compute the max flow?

What is a 'network flow'?

- Directed graph, where each edge has a max capacity
- Amount of 'flow' into the node = amount of 'flow' out of the node, unless it's the source or sink
- Used for modelling traffic systems, fluids in pipes, bits across wires...
- Can we indeed model the Software Allocation problem as a network flow one? Show it.
- Then: which algorithm to compute the max flow?
- Edmonds-Karp's algorithm <https://www.youtube.com/watch?v=MczX0SM3I84>
- See also §4.6 of CP3 for a description

Outline

- 1 Graphs
 - As long as I learn, I live
 - Railway construction
 - Software Allocation
- 2 Data structures
 - Streaming System
 - Young and Successful
 - Hic-Hac-Hoe
- 3 Strategy and teamwork

Streaming System

- This is UVa problem 1203, “Argus” (and also used in the Asia - Beijing - 2004/2005 ACM ICPC Local Contest). (ICPC live 6421)
- Main point of this exercise here: Data structures
- Solve it

Streaming System

- This is UVA problem 1203, “Argus” (and also used in the Asia - Beijing - 2004/2005 ACM ICPC Local Contest). (ICPC live 6421)
 - Main point of this exercise here: Data structures
 - Solve it
- ⇒ Priority Queue

Young and Successful

- Not an ACM ICPC problem from recent years (or at all)
- Description on printout
- Identify type
- Try to solve it

Young and Successful

- Not an ACM ICPC problem from recent years (or at all)
- Description on printout
- Identify type
- Try to solve it
- Solve it algorithmically

Young and Successful

- Not an ACM ICPC problem from recent years (or at all)
 - Description on printout
 - Identify type
 - Try to solve it
 - Solve it algorithmically
- ⇒ Basically a version of a *range search*:
input: A set S with n points in d -dimensional space E , and a query asking for the points in region Q

Hic-Hac-Hoe

- Slight modification to Tic-Tac-Toe: infinite board size
- data structure & search, with a time limit of 2 sec.
- Solve it

Hic-Hac-Hoe

- Slight modification to Tic-Tac-Toe: infinite board size
 - data structure & search, with a time limit of 2 sec.
 - Solve it
- ⇒ instead of a 2D array, use a **balanced BST** to store the o's and x's, and use that BST to check the game state
- ⇒ Balanced BST amounts to using the Divide & Conquer strategy
- Note: difference with heap is that the tree doesn't have to be complete
 - C++ STL: map and set; Java with TreeMap and TreeSet

Outline

- 1 Graphs
 - As long as I learn, I live
 - Railway construction
 - Software Allocation
- 2 Data structures
 - Streaming System
 - Young and Successful
 - Hic-Hac-Hoe
- 3 Strategy and teamwork

Looking at a whole contest

- The 2013 ACM-ICPC Arab Regional Programming Contest
- Groups of 3. Consider:
 - Roles: who will do what? Manager/time-keeper?
 - Pair solving and pair programming
 - Categorise problems (recall problem solving strategies)
 - Know your strengths and weaknesses
 - Decide which problem(s) to solve first
- <https://icpcarchive.ecs.baylor.edu/>, “Browse Problems” – “Regionals 2013” – “Africa/Middle East” etc.

Scheduled training dates

- Aug 6: 10:00-16:00
- Aug 13: 10:00-16:00
- Aug 27: 10:00-15:00
- **Sept 10: 10:00-16:00** → **Ashraf Moolla on DP**
- Sept 17: 10:00-16:00
- Sept 24: 10:00-16:00 or Oct 1: 10:00-16:00
- Date of the regionals: very likely Oct 15